Vampir

MPI-Performance Analysis

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How does one gain insight about an HPC applciation?

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Data has to be presented to the developer in an understandable way in order to gain insight!

- **Data can be broken up**
  - by time
    - Example: Function entries and exits are displayed on a timeline
  - attributed to graphs representing properties of the source code
    - Example: Display function execution times on the call graph
  - etc…

Introduction to Vampir

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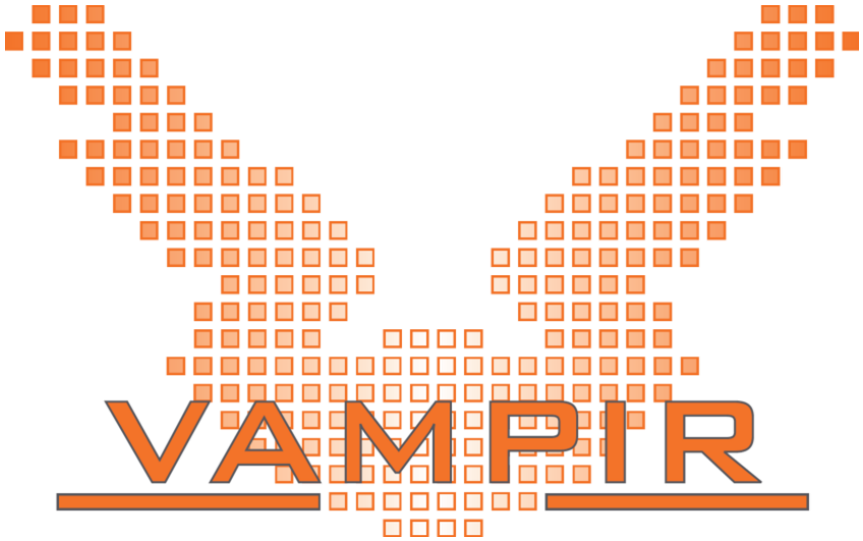
Overview of Vampir

How to use Vampir



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Facts about Vampir

Actually a tool collection

- Vampir Trace — freely available
- Vampir GUI — commercial

Developed at the TU Dresden

Many additional recording capabilities

- Memory Usage
- I/O Activity Tracing
- CPU-counters

Prepare your environment

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Vampir-trace

```
module load vampirtrace
```

Vampir (GUI)

```
module load vampir
```

Note:

- **Vampir-Trace is not necessary for visualization of data**
- **The Vampir module must be loaded for analysis**

How to use

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To use Vampir the following steps are necessary:

1. **Select the correct compiler wrapper**

C:
```
vtcc  -vt:cc
```

C++:
```
vtcxx  -vt:cxx
```

Fortran:
```
vtf90  -vt:f90
```
```
vtf77  -vt:f77
```

2. **Instrument your application by recompiling it**

```
vtcc –vt:cc gcc hello.c –o hello.exe
```

How to use

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3. Execute as usual

```
$MPIEXEC  -np 2 hello.exe
```

⇨ Results in a OTF-file

4. Analyze application

```
vampir vtrace.otf
```

Interlude #1: Measurement Basics

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Two independent decisions:

1. When performance is measurement triggered

- Sampling
  - Triggered by timer interrupt or by hardware counter overflow
  - Can measure unmodified executables, potential low overhead
- Code instrumentation:
  - Triggered by “instrumentation hooks” inserted into the code
  - Insertion can be done manually or automatically

2. How performance is data recorded

- Profile ::= Summarization of events over time
  - run time summarization (functions, call sites, loops, …)
- Trace file ::= Sequence of events over time

A study of a short example program

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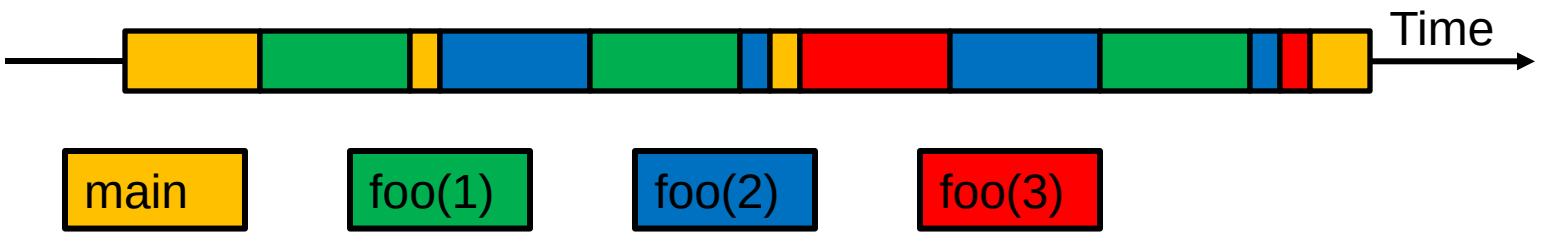
To explain the differences:

```
main(..)
{
  for (i=1..3)
  {
    foo(i)
  }
}

foo(i)
{
  if (i>0)  foo(i-1)
}
```



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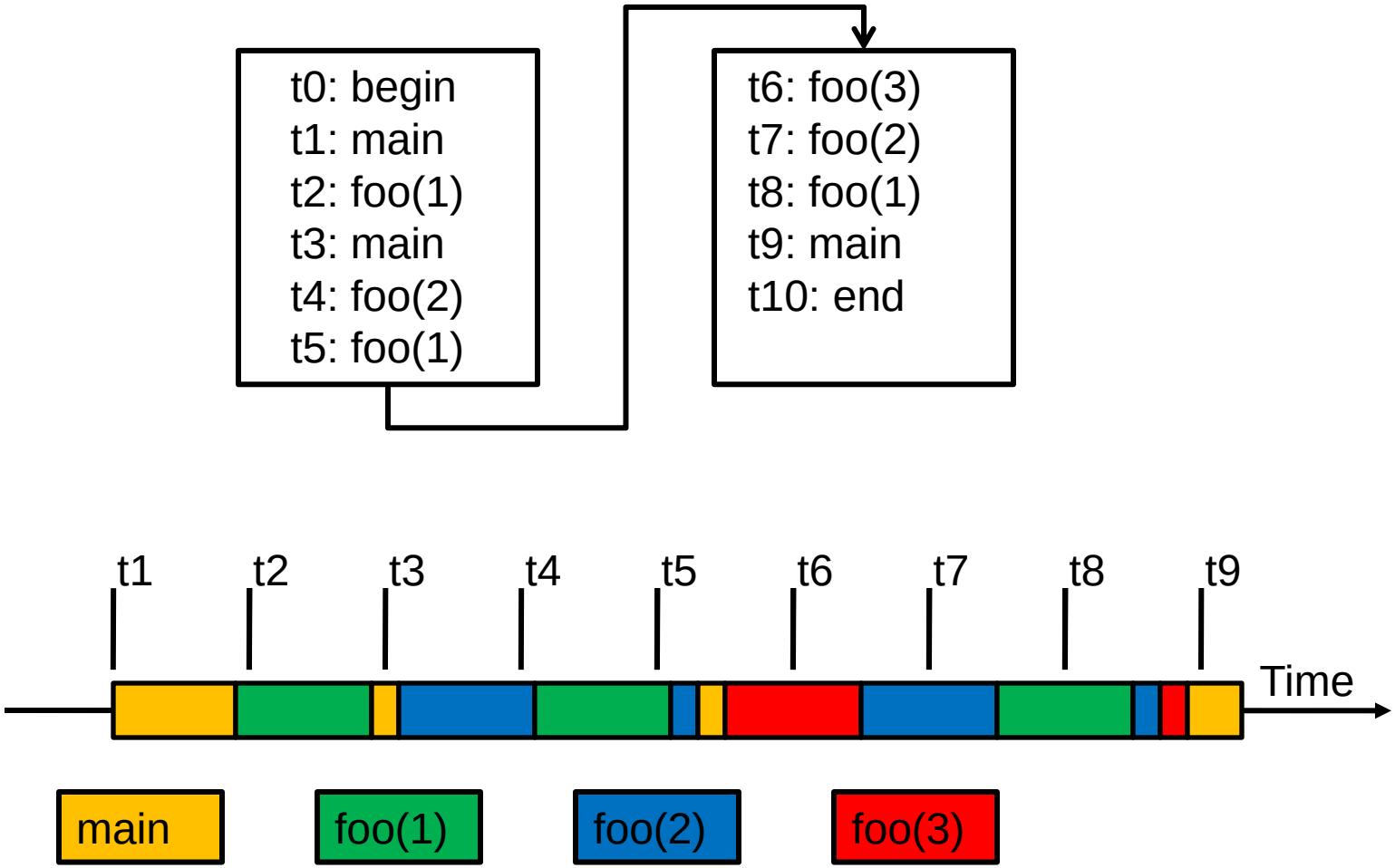
Example: Sampling



Summary: Sampling

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Sampling:

The application is probed at specific times and a set of interesting metrics is gathered

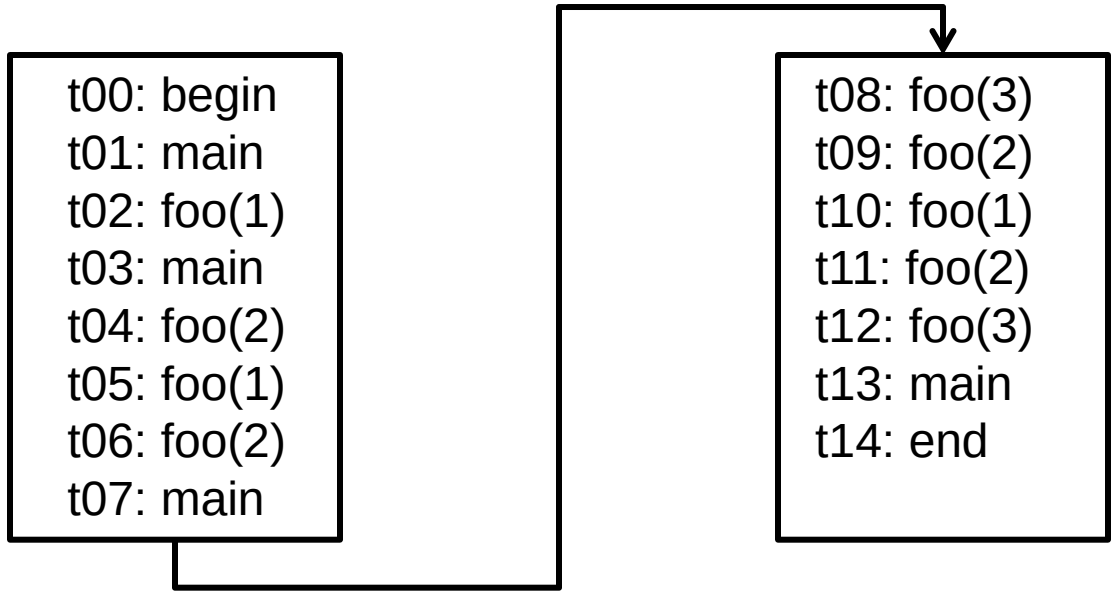
Advantages:

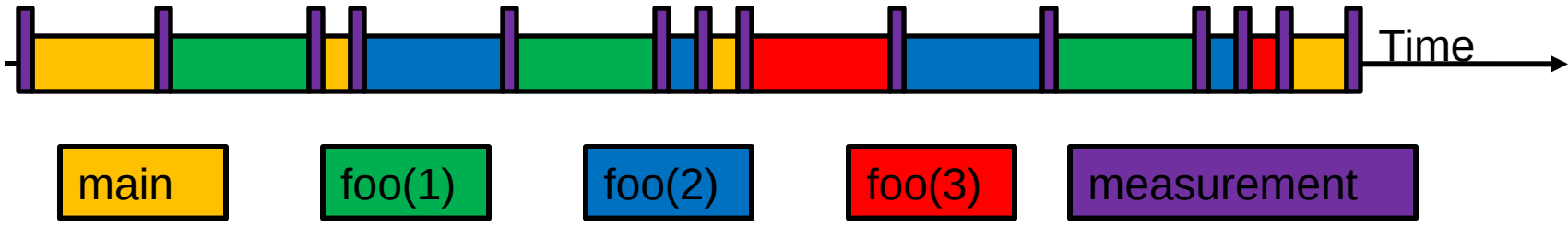
+ **Low perturbance of the application**
+ **Application does not have to be recompiled**
+ **Works well with large and long running applications**

Disadvantages:

- **Not very detailed information on high frequency metrics**

Example: Instrumentation





Summary: Instrumentation

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Idea:

The code is instrumented such that every interesting event is recorded as it occurs.

Advantages:

+ **Every event of interest can be captured**
+ **Much more detailed information possible**

Disadvantages:

- **Preprocessing of the program necessary**
- **Probably expensive at runtime**

Critical issues

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Accuracy

- Perturbation

  Measurement alters program behaviour

- Intrusion overhead

  Measurement itself needs time and thus lowers performance

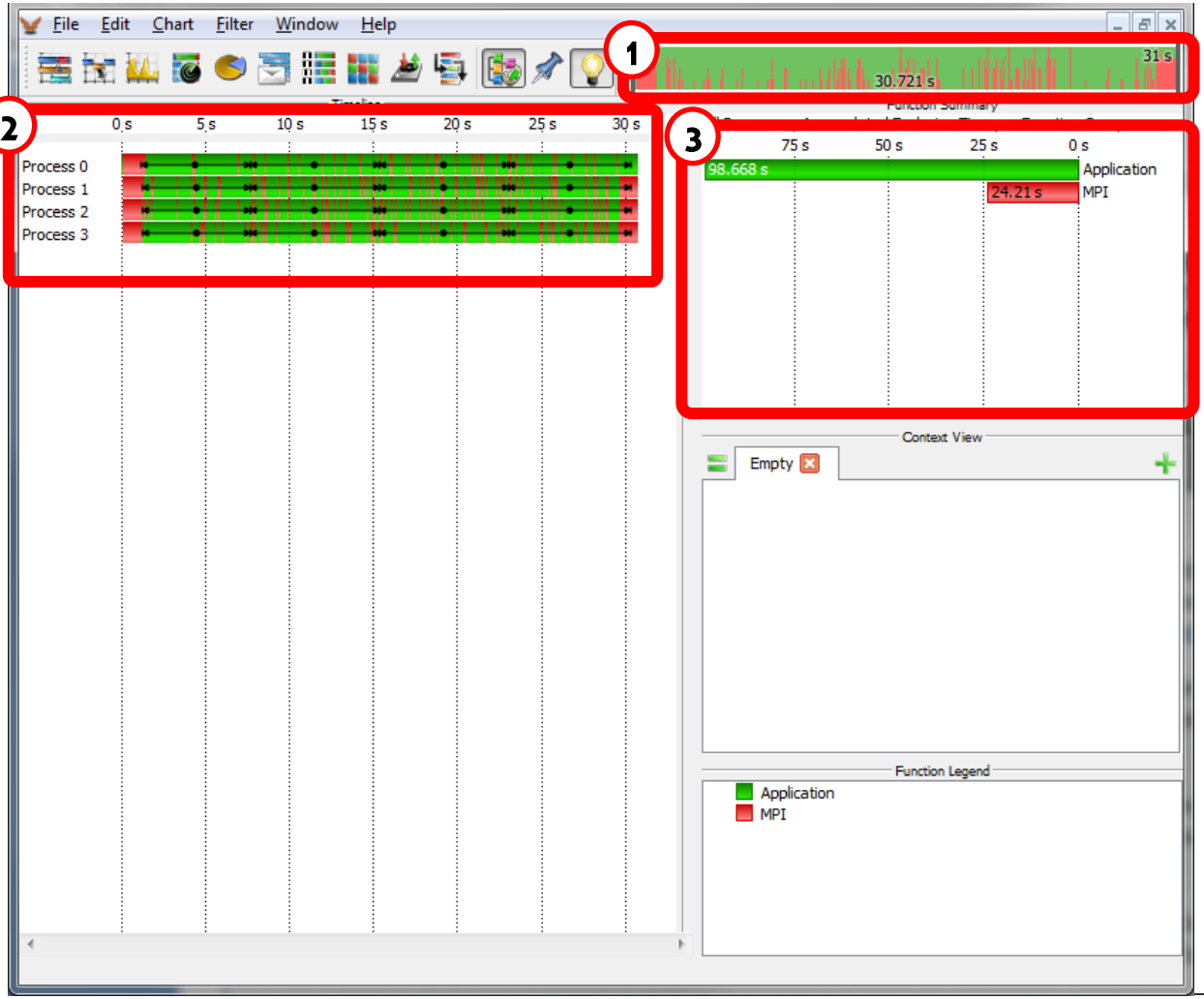
- Accuracy of timers, counters

Granularity

- How many measurements
- How much information / work during each measurement

Trade-off: **Accuracy ⇔ Expressiveness of data**

Vampir - Main Window



1) **Overview Timeline**
2) **„Master“ Timline**
3) **Function Summary**

Function Summary:
- all data is grouped
- aggregated for all processes

To Access a pannels options: **Right Click in the Window**

Important Vampir Charts

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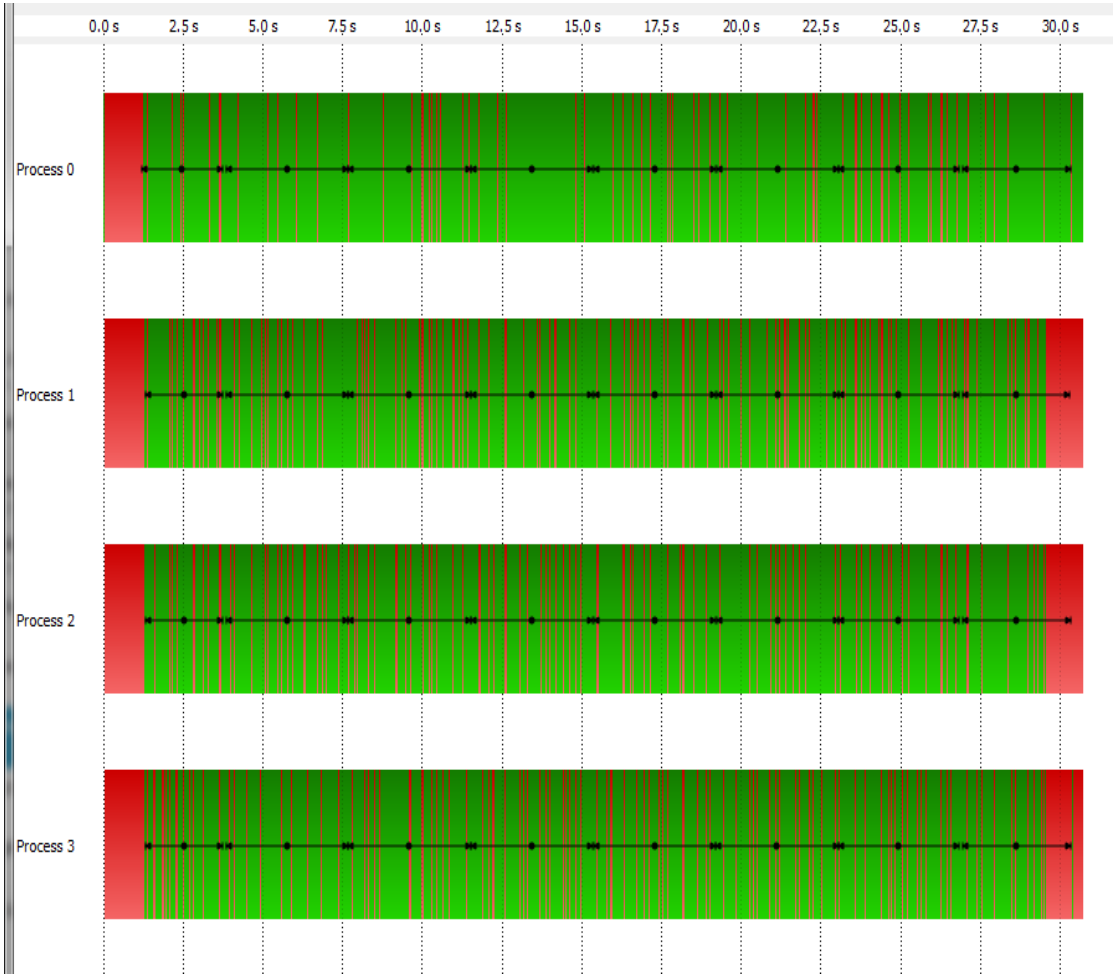
Main Charts

1) Master Timeline
2) Process Timeline
3) Counter Timeline
4) Function Summary
5) Message Summary
6) Process Summary
7) Communication Matrix
8) Call Tree

Master Timeline



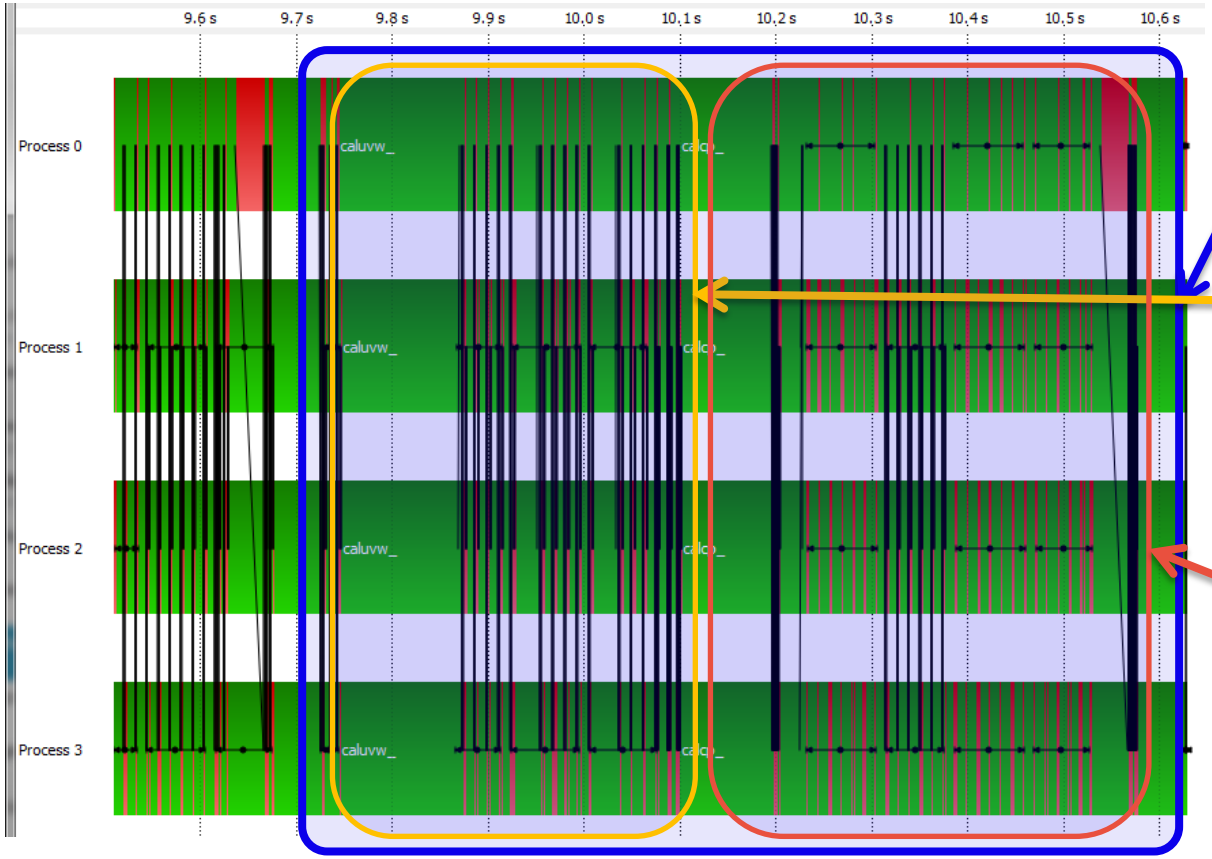
- **Timeline for every process (and thread)**

Possible Options:

- **Vertical Zooming**
- **Horizontal Zooming**

- **Vertical / Horizontal Panning**

Master Timeline (more detail)

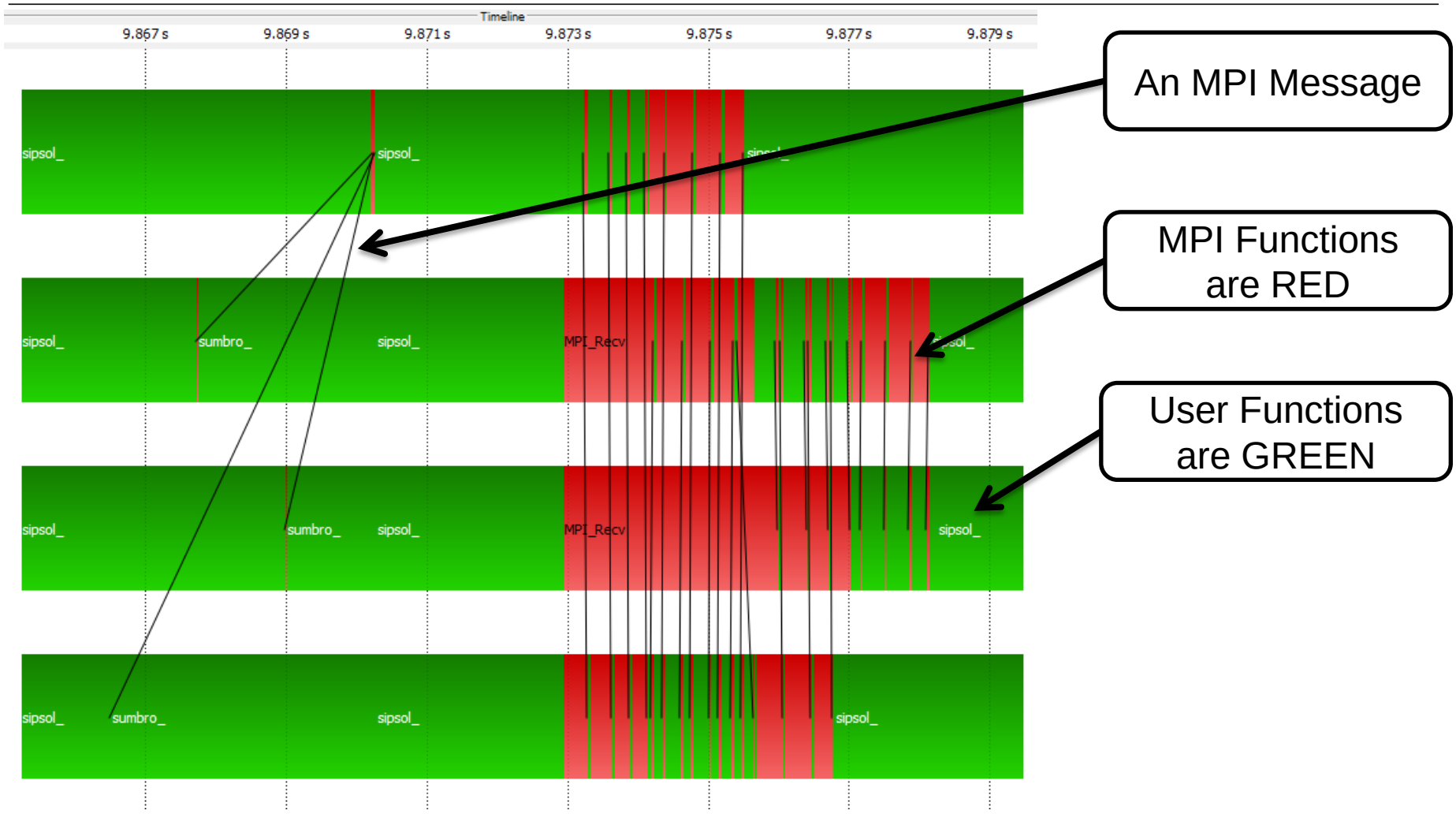


Single FASTEST Iteration

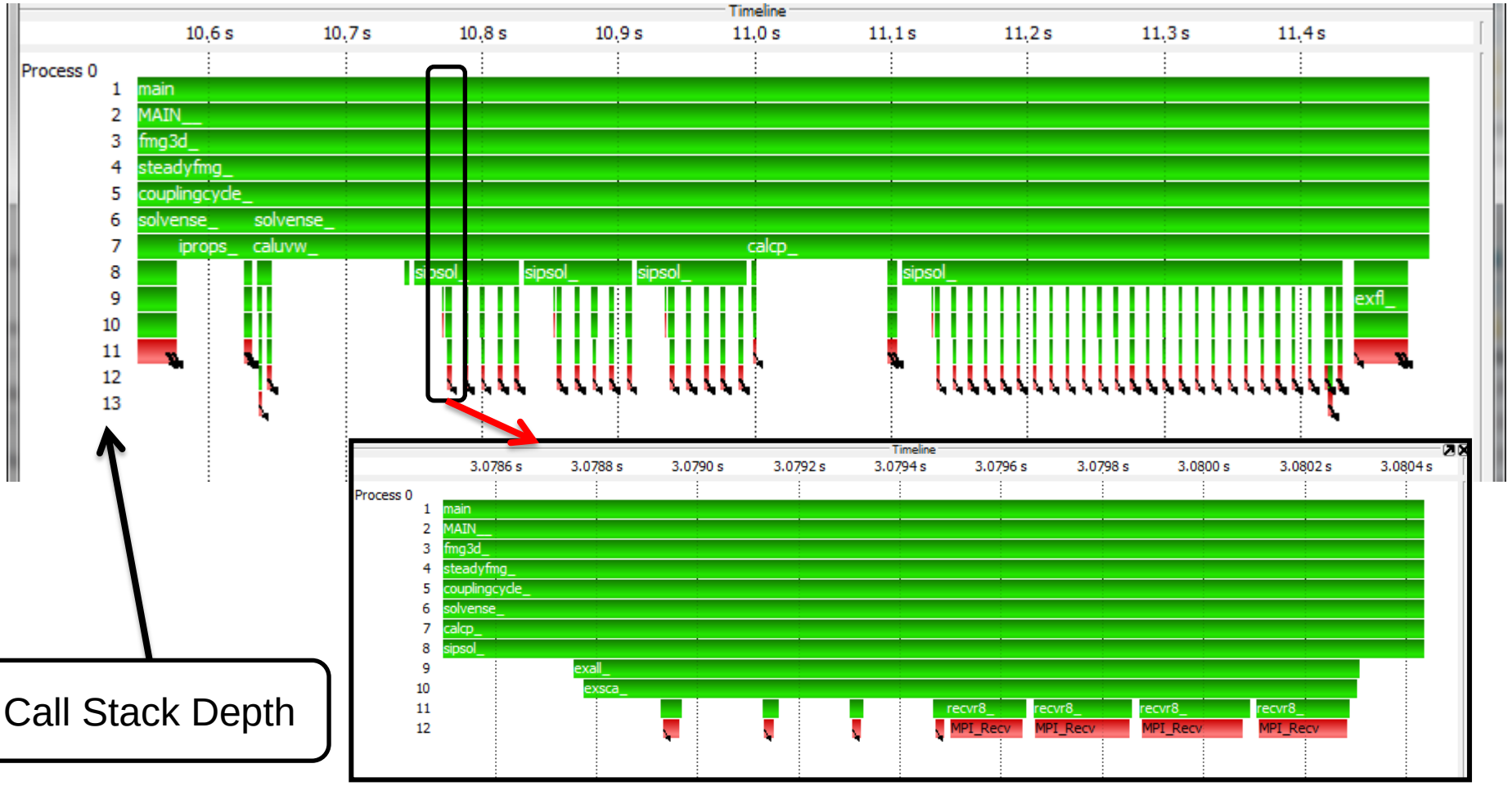
Solving of Mass Transport/ Flow Equation

Solving of Mass Transport/ Flow Equation

Master Timeline (even more detail)



Process Timeline



Call Stack Depth

Interlude #2:
What are Hardware Performance Counter?

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- **Hardware Performance Counter (Wikipedia):**

  … a set of (programmable) special-purpose registers built into modern microprocessors to store the counts of hardware-related activities within computer systems …

- **What do they do?**

  Example MEGAFLOPS

  - MEGA = Million (10^6)
  - FLOPS = FLoatingpoint Operations Per Second
  - Tell the CPU to count the number of floating point operations

Excerpt of available counters Nehalem X5570

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- Level 1 data cache misses
- Level 1 instruction cache misses
- Level 2 data cache misses
- Level 2 instruction cache misses
- Level 3 data cache misses
- Level 3 instruction cache misses
- Level 1 cache misses
- Level 2 cache misses
- Level 3 cache misses
- Requests for a snoop
- Requests for exclusive access to shared cache line
- Requests for exclusive access to clean cache line
- Requests for cache line invalidation
- Requests for cache line intervention
- Level 3 load misses
- Level 3 store misses
- Cycles branch units are idle
- Cycles integer units are idle
- Cycles floating point units are idle
- Cycles load/store units are idle
- Data translation lookaside buffer misses
- Instruction translation lookaside buffer misses
- Total translation lookaside buffer misses
- Level 1 store misses
- Level 1 load misses
- Level 2 load misses
- Level 2 store misses
- Branch target address cache misses
- Data prefetch cache misses
- Level 3 data cache hits
- Translation lookaside buffer shootdowns
- Failed store conditional instructions
- Successful store conditional instructions
- Total store conditional instructions
- Cycles Stalled Waiting for memory accesses
- Cycles Stalled Waiting for memory Reads
- Cycles Stalled Waiting for memory writes
- Cycles with no instruction issue
- Cycles with maximum instruction issue
- Cycles with no instructions completed
- Cycles with maximum instructions completed
- Hardware interrupts
- Unconditional branch instructions
- Conditional branch instructions
- Conditional branch instructions taken
- Conditional branch instructions not taken
- Conditional branch instructions mispredicted
- Conditional branch instructions correctly predicted
- FMA instructions completed

Excerpt of available counters Nehalem X5570

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- Instructions issued
- Instructions completed
- Integer instructions
- Floating point instructions
- Load instructions
- Store instructions
- Branch instructions
- Vector/SIMD instructions (could include integer)
- Cycles stalled on any resource
- Cycles the FP unit(s) are stalled
- Total cycles
- Load/store instructions completed
- Synchronization instructions completed
- Level 1 data cache hits
- Level 2 data cache hits
- Level 1 data cache accesses
- Level 2 data cache accesses
- Level 3 data cache accesses
- Level 1 data cache reads
- Level 2 data cache reads
- Level 3 data cache reads
- Level 1 data cache writes
- Level 2 data cache writes
- Level 3 data cache writes
- Level 1 instruction cache hits
- Level 2 instruction cache hits
- Level 3 instruction cache hits
- Level 1 instruction cache accesses
- Level 2 instruction cache accesses
- Level 3 instruction cache accesses
- Level 1 instruction cache reads
- Level 2 instruction cache reads
- Level 3 instruction cache reads
- Level 1 instruction cache writes
- Level 2 instruction cache writes
- Level 3 instruction cache writes
- Level 1 total cache hits
- Level 2 total cache hits
- Level 3 total cache hits
- Level 1 total cache accesses
- Level 2 total cache accesses
- Level 3 total cache accesses
- Level 1 total cache reads
- Level 2 total cache reads
- Level 3 total cache reads
- Level 1 total cache writes
- Level 2 total cache writes
- Level 3 total cache writes
- Floating point multiply instructions
- Floating point add instructions
- Floating point divide instructions

Excerpt of available counters Nehalem X5570

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- Floating point square root instructions
- Floating point inverse instructions
- Floating point operations
- Floating point operations; optimized to count scaled single precision vector operations
- Floating point operations; optimized to count scaled double precision vector operations
- Single precision vector/SIMD instructions
- Double precision vector/SIMD instructions

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Counter Timeline

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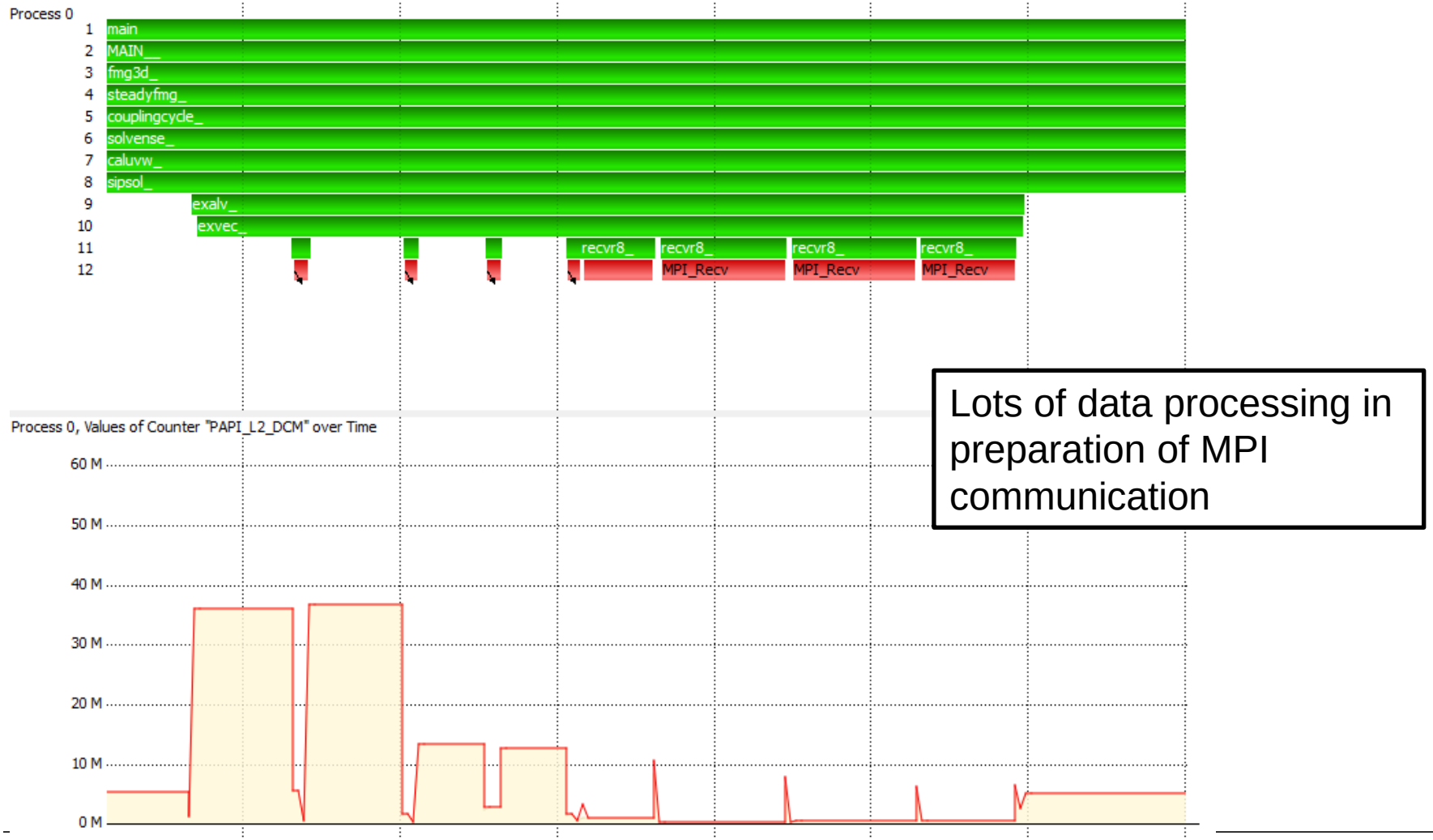

Lots of data processing in preparation of MPI communication

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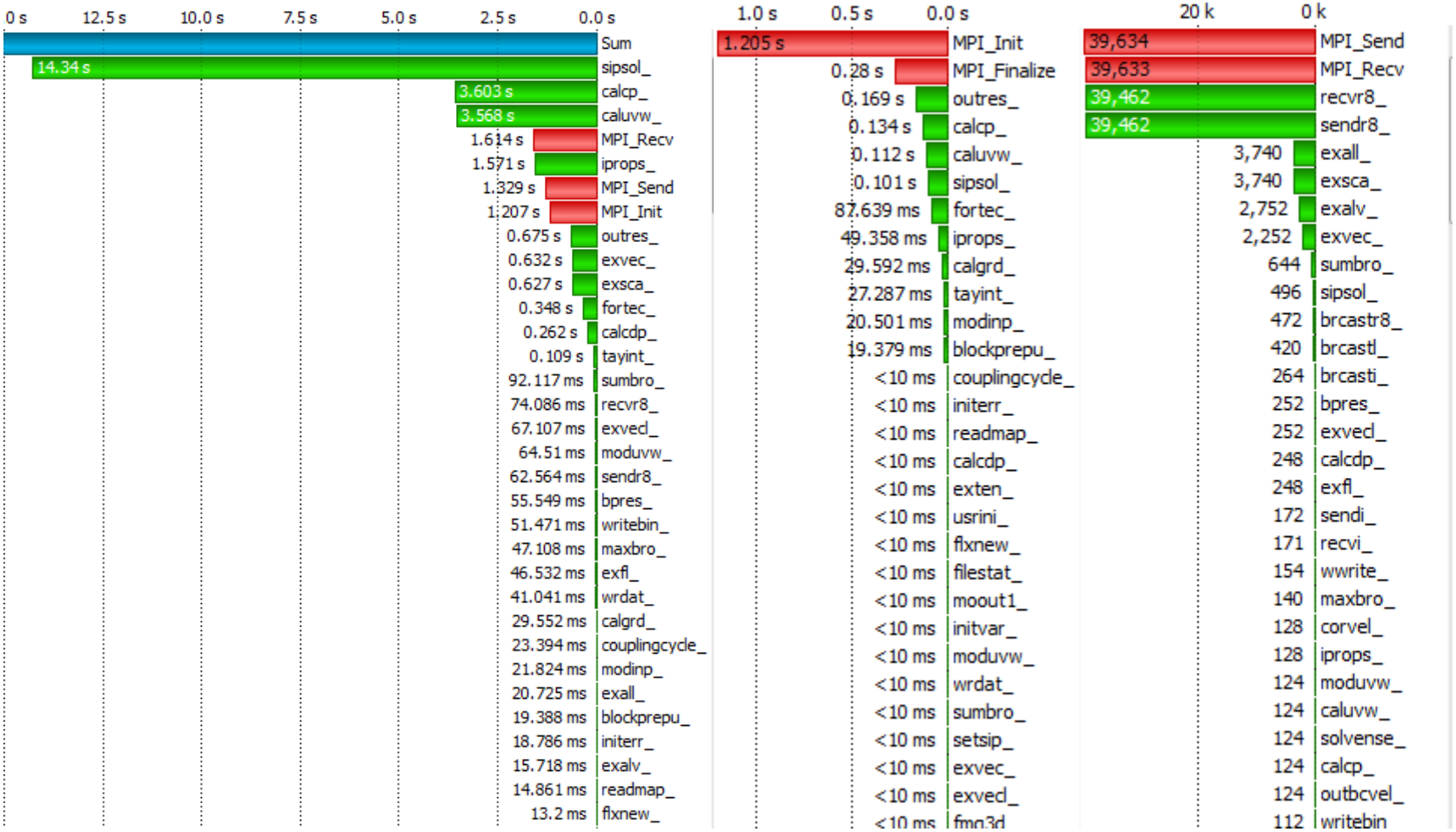
Function Summary



Message Summary & Process Summary

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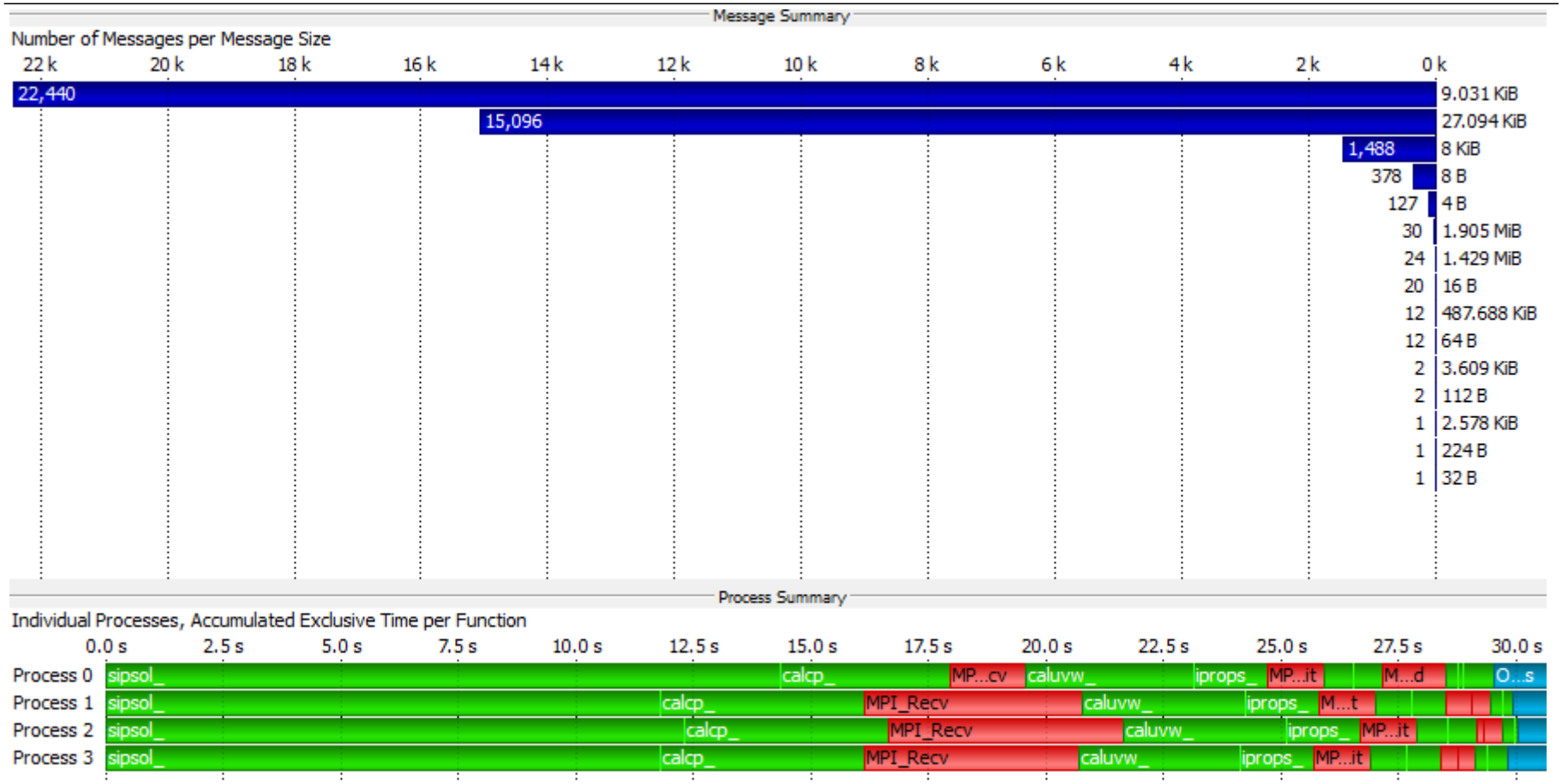

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Call Tree

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Call Tree

All Processes

Function	Min Number of Invocations	Max Number of Invocations	Min Inclusive Time	Max Inclusive Time
gridoutput_	1	1	0.785 s	1.176 s
flxin_	1	1	219.593 µs	755.892 µs
couplingcycle_	1	1	27.953 s	27.953 s
wwrite_	35	35	146.162 µs	3.106 ms
solvense_	31	31	27.876 s	27.925 s
outbcvel_	31	31	7.300 ms	89.991 ms
iprops_	31	31	1.675 s	1.683 s
corvel_	31	31	235.623 µs	246.170 µs
caluvw_	31	31	11.102 s	11.187 s
calcp_	31	31	14.999 s	15.053 s
brcastl_	31	31	191.570 µs	281.082 µs
reread_	1	1	40.274 µs	121.047 µs
maxbro_	31	31	544.538 µs	47.062 ms
gtdate_	0	2	0.000 s	74.938 µs
checkprop_	1	1	58.620 µs	59.728 µs
brcastl_	31	31	324.689 µs	22.578 ms
bpres_	1	1	3.026 ms	3.150 ms
initialize_	1	1	1.583 s	1.587 s
setvec_	1	1	894.130 µs	948.701 µs
setsip_	1	1	972.842 µs	1.603 ms
setprefindx_	1	1	1.783 ms	7.207 ms
readprop_	1	1	529.421 µs	551.436 µs
brcastr8_	10	10	59.762 µs	487.891 µs
brcasti_	1	1	6.422 µs	8.351 µs
readmap_	1	1	32.387 ms	32.449 ms
sendr8_	0	9	0.000 s	16.857 ms
sendi_	0	39	0.000 s	619.185 µs
recvr8_	0	3	0.000 s	5.672 ms
recvi_	0	13	0.000 s	22.103 ms
brcastl_	1	1	5.907 µs	8.126 µs

Vampir also provides runtime filtering capabilities

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1. Define a filter file in the environment variable

   export VT_FILTER_SPEC=~/myFilter

2. The filter file contains a list of functions names with a limit
3. Execute your application as usual

```
yy_get_next_buffer -- 0
yy_get_previous_state -- 0
ewdlocalize_ -- 5000
ewdlocalizei_ -- 0
ewdfeshape_ -- 67000
ewdfeshapest3d4n_ -- 0
yylex -- 0
ewdzero_ -- 0
ewdbsradd_ -- 0
ewdbsraddi_ -- 0
```